

Shapely monads for graphical calculi

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Outline

- ① Introduction
- ② Preliminaries
- ③ Operads
- ④ Graphical operads
- ⑤ Shapely monads

Mathematical motivation

Certain algebraic structures with

- obvious graphical intuition;
- tedious formal definition.

E.g., operads, properads, polycategories, PROPs, and variants.

Computer science motivation

Graphical calculi with

- obvious graphical intuition;
- tedious formal definition;
- involved or non-existent notion of model.

E.g., took quite long to work out for proof nets¹!

Example

Interaction nets, (multiplicative) proof nets, bigraphs, ZX-calculus.

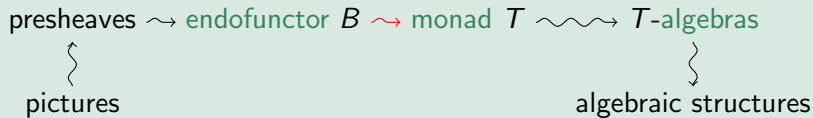
¹Bierman. *On Intuitionistic Linear Logic*. PhD thesis, Cambridge, 1993.

Contributions

- Make graphical intuition rigorous thanks to **presheaf** theory.
- \rightsquigarrow Alternative definition of
 - maths**: the algebraic structure in question
 - comp. sci.**: a notion of model for the graphical calculus in question.
- View old definition as economical characterisation:

	old definition	new definition
statement	hard	easy
construction	easy	hard

Posing the problem categorically



Need to explain these terms, at least intuitively.

- Rightmost part: standard categorical approach to algebra.
- Difficult part in red!

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Categories

Definition

Objects, and morphisms between them.

Example

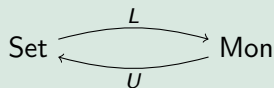
	Objects	Morphisms
Set	Sets	Functions
Mon	Monoids	Monoid homomorphisms
Grp	Groups	Group homomorphisms
...		

Functors

Definition

Functor = morphism of categories.

Example



- Action on objects:

$$L(X) = \sum_n X^n$$

= sequences of elements of X ,
= **free monoid** on X .

Multiplication:

$$(x_1, \dots, x_n), (x_{n+1}, \dots, x_p) \mapsto (x_1, \dots, x_p).$$

- Action on morphisms:

$$\begin{aligned} L(X \xrightarrow{f} Y) : L(X) &\rightarrow L(Y) \\ (x_1, \dots, x_n) &\mapsto (f(x_1), \dots, f(x_n)). \end{aligned}$$

- Other example:

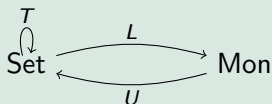
$$U(M) = |M|, \text{ carrier of } M.$$

Monads

Definition

Monad = endofunctor + structure.

Example



- Composite $T = U \circ L$.
- $T(X)$ = free monoid viewed as a set.
- T is a monad.

Crucial point I: algebraic structures = algebras for a monad

T -algebra

T -algebra = morphism
$$\begin{array}{c} T(X) \\ m \downarrow \\ X \end{array}$$
 with easy conditions.

Example: previous T

- $T(X)$ = free monoid viewed as a set.
- So m maps sequences (x_1, \dots, x_n) to elements.
- Thought of as multiplication.

Example T -algebra:
$$\begin{array}{l} m: T(\mathbb{N}) \rightarrow \mathbb{N} \\ (n_1, \dots, n_p) \mapsto \sum_i n_i. \end{array}$$

Morphisms of algebras

Morphisms of T -algebras

$$\begin{array}{ccc}
 T(X) & \xrightarrow{T(f)} & T(Y) \\
 m \downarrow & & \downarrow m' \\
 X & \xrightarrow{f} & Y
 \end{array}$$

- $f(m(x_1, \dots, x_n)) = m'(f(x_1), \dots, f(x_n))$.
- Morphism = structure-preserving map.

Proposition (in the monoids example)

T -algebras form a category $T\text{-Alg}$, **equivalent** to Mon .

Moral (standard, but very important!)

Algebraic structure (monoids) \Leftarrow monad T .

T describes 'free' algebraic structures.

Other examples on sets

Algebraic structure	$T(X)$
Monoids	$\sum_n X^n$
Commutative monoids	$\sum_n X^n / \mathfrak{S}_n$
Rings, modules, algebras,
Complete semi-lattices	$\mathcal{P}(X)$

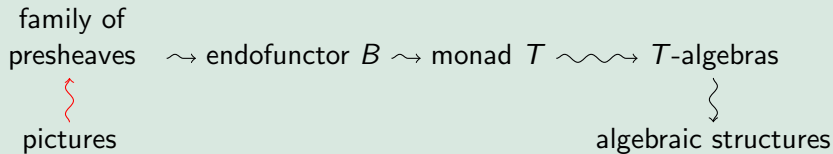
Non-example: fields, as there are no free fields over a set.

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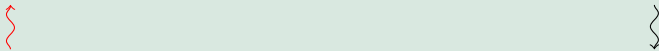
From pictures to presheaves

- Running example: (nonsymmetric, coloured) operads.
- Well-known case: T already known!
- Result specialises to: characterisation of T as a free **shapely** monad.



From pictures to presheaves

- Running example: (nonsymmetric, coloured) operads.
- Well-known case: T already known!
- Result specialises to: characterisation of T as a free **shapely** monad.

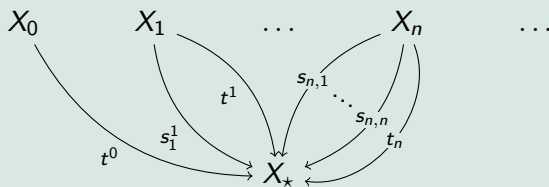
family of
 multigraphs \rightsquigarrow endofunctor $B \rightsquigarrow$ monad $T \rightsquigarrow T$ -algebras


 pictures algebraic structures

Multigraphs

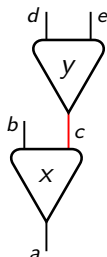
Multigraph $X \approx$ graph whose edges may have several sources.

Diagram



- X_* : vertices;
- X_n : edges with n sources;
- $s_{n,i}(e)$: i th source of n -ary e ;
- $t_n(e)$: target of e .

Example multigraph



- $X_* = \{a, b, c, d, e\}$,
- $X_2 = \{x, y\}$,
- $X_n = \emptyset$ otherwise,
- $t_2(x) = x \cdot t = a$ (notation!),
- $x \cdot s_1 = b$, $x \cdot s_2 = c$, $y \cdot t = c$,
 $y \cdot s_1 = d$, $y \cdot s_2 = e$.

Category of multigraphs

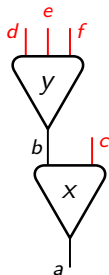
Morphism = map preserving target and sources.

Proposition

Multigraphs form a category MGph .

Intuitive definition

A (*nonsymmetric, coloured*) operad (in sets) \mathcal{O} is a multigraph \mathcal{O} with 'plugging', e.g., for all $x \in \mathcal{O}_2$ and $y \in \mathcal{O}_3$ with $y \cdot t = x \cdot s_1$, one may form



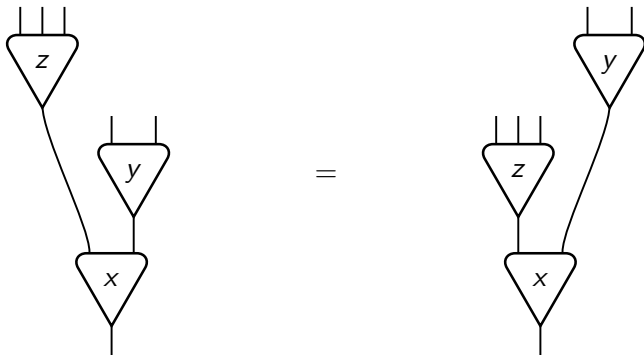
in \mathcal{O}_4 .

Notation

Denoted by $x \circ_1^{2,3} y$.

Intuitive definition (cont'd)

Plugging should satisfy obvious graphical axioms, e.g.,



Dreadful glimpses of standard definition

Definition

A (*nonsymmetric, coloured*) operad (in sets) is

- a multigraph \mathcal{O} , together with
- for all m, n, i , $x \in \mathcal{O}_m$ and $y \in \mathcal{O}_n$ such that $x \cdot s_i = y \cdot t$, an element

$$x \circ_i^{m,n} y \in \mathcal{O}_{m+n-1};$$

- for all $a \in \mathcal{O}_*$, an element $id_a \in \mathcal{O}_1$;
- satisfying axioms like

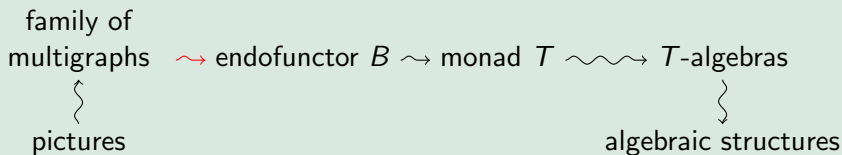
$$(x \circ_i^{m,n} y) \circ_j^{m+n-1,p} z = \begin{cases} (x \circ_j^{m,p} z) \circ_{i+p-1}^{m+p-1,n} y & (\text{if } j < i) \\ x \circ_i^{m,n+p-1} (y \circ_{j-i+1}^{n,p} z) & (\text{if } i \leq j < i+n) \end{cases}$$

for all $x \in \mathcal{O}_m$, $y \in \mathcal{O}_n$, $z \in \mathcal{O}_p$.

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Endofunctors from multigraphs



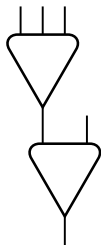
Crucial point II: arguments for composition = multigraph morphisms

- Recall the picture for composition in \mathcal{O} , on the right.
- View it as a **multigraph**, say X .

(Morphisms $X \rightarrow \mathcal{O}$) \Leftrightarrow (choices of (x, y)):

- $x \in \mathcal{O}_2$ and $y \in \mathcal{O}_3$,
- such that $x \cdot s_1 = y \cdot t$.

= potential arguments for $\circ_1^{2,3}$ if it existed.



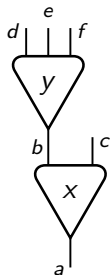
Arities

Definition (Basic arities)

- X is the **arity** of $\circ_1^{2,3}$.
- Obvious generalisation: $X_i^{m,n}$ is the **arity** of $\circ_i^{m,n}$.
- Similarly, arity of id : multigraph with just one vertex (wire).

Making sense of h_X -algebras

- Recall our example multigraph X on the right.
- Consider the functor $h_X : \text{MGph} \rightarrow \text{MGph}$ defined by:
 - $h_X(Y)_* = Y_*$,
 - $h_X(Y)_4 = \text{MGph}(X, Y)$, the set of multigraph morphisms from X to Y ,
 - $h_X(Y)_n = \emptyset$ for $n \neq 4$.
- So $h_X(Y)_4 = \{(x', y') \in Y_2 \times Y_3 \mid x' \cdot s_1 = y' \cdot t\}$.
- An algebra $h_X(Y) \rightarrow Y$ is determined by:
 - a multigraph Y ,
 - plus a map $h_X(Y)_4 \rightarrow Y_4$, i.e.,
 - an interpretation of $\circ_1^{2,3}$!

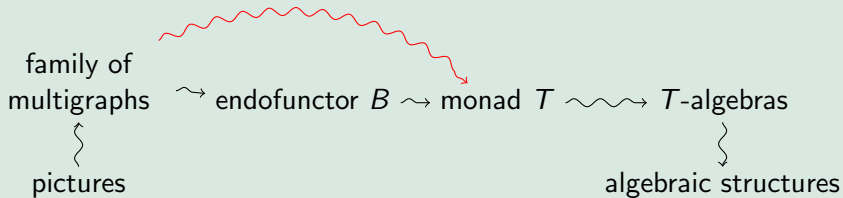


Summary

Multigraph $X \rightsquigarrow$ functor which specifies an operation of arity X .

I.e., algebras have such an operation.

The monad from derived arities



Graphical definition of operads

Need to define arities for all **derived** operations:

Definition

Let \mathcal{T}_n denote the class of planar trees with n leaves.

Define $T: \text{MGph} \rightarrow \text{MGph}$ by:

- $T(Y)_\star = Y_\star$,
- $T(Y)_n = \sum_{X \in \mathcal{T}_n} \text{MGph}(X, Y)$, the set of multigraph morphisms from **some** n -ary tree X to Y .

Lemma

The functor T is a monad on MGph .

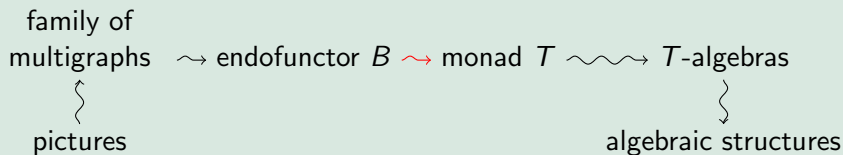
Theorem

Operads are equivalent to T -algebras.

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Generating monads



- Goal: generate T automatically from **basic arities**.

- ▶ Compositions $X_i^{n,m}$.
- ▶ Identities I_a .

Signature for operads

Definition

Let \mathcal{B}_n denote the set of basic arities with n leaves.

Intuition: filiform trees of depth 2.

Define $B: \text{MGph} \rightarrow \text{MGph}$ by:

- $B(Y)_\star = Y_\star$,
- $B(Y)_n = \sum_{X \in \mathcal{B}_n} \text{MGph}(X, Y)$, the set of multigraph morphisms from some n -ary basic arity X to Y .

Question: how to generate T from B ?

Naive attempt

Well-known correspondence

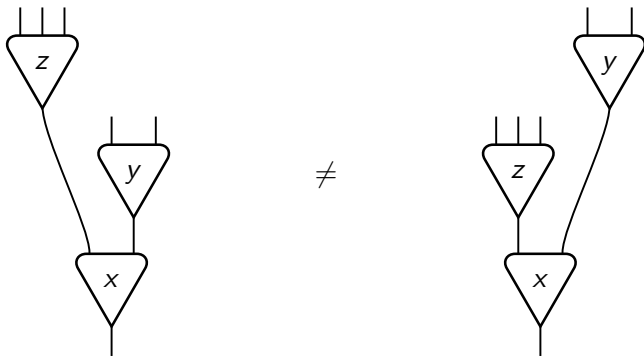


Miss!

$$\mathcal{M}(B) \not\cong T.$$

Reason

$\mathcal{M}(B)$ -algebras do not satisfy any of the axioms!



Which monads **do** enforce them? **Shapely** ones!

Shapely monads

Subcategory

$$\text{Framed}(\text{MGph}) \subseteq \text{Cell}(\text{MGph}) \subseteq \text{Analytic}(\text{MGph}) \subseteq \text{Endo}(\text{MGph}).$$

- Stable under composition.
- Has a terminal object \top , automatically a monad.

Definition

Shapely = subfunctor of \top in $\text{Framed}(\text{MGph})$.

Graphical calculus = shapely monad.

Intuition: at most one operation of each arity.

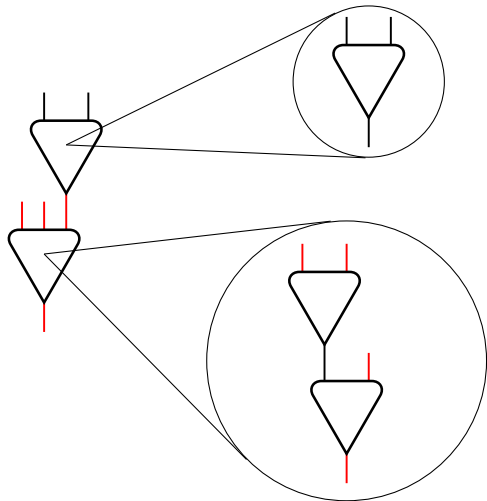
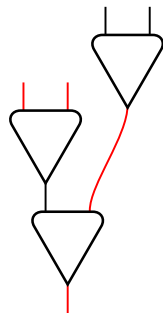
Generation result

Theorem

$T = \bigcup_n (id \cup B)^{\cdot n}$ is the free shapely monad over B .

$B \cdot B$ denotes the **image** of $B \circ B$: $B \circ B \twoheadrightarrow B \cdot B \hookrightarrow T$.

Illustration of $B \cdot B$

 $B \circ B$

 \rightarrow
 $B \cdot B$
 \mapsto


General result

- Consider any presheaf category with a **subterminal** object \top .

At most one morphism from any object to \top .

- Consider \top -**analytic** functors, i.e., analytic functors with a map to \top .
- Suppose they are stable under composition.
- Example: framed endofunctors.

Definition

Shapely functor = subfunctor of \top .

Theorem

The free shapely monad on a shapely endofunctor B is $\bigcup_n (id \cup B)^{\cdot n}$.

Applications

- Characterisation of the monads for polycategories, properads, PROPs, etc, as free shapely monads.
- Definition of free shapely monads for interation nets and fragments of proof nets.

Conclusion

If you wonder what a model for your graphical calculus is, you could try to:

- Formalise pictures as presheaves.
- Derive monad automatically from them.
- Then work on an intelligible characterisation of algebras.

Perspectives

- Further applications, e.g., to PROP rewriting (started long ago with Adrien D., idle).
- Framed functors: 2-levels, level 1 fixed. Generalisation?
- Notion of **representable** algebra, as in representable operad.

Existence of a tensor product.

- Notion of **weak** algebra (suggested by Kris W.), as in ∞ -operad.
- Generalise to not strictly shapely, e.g., proof net boxes.

Thanks!



Shapely functors: intuition

- Restrict to functors with at most one operation per arity.
- There should be one ‘full’ such functor \mathbb{T} , with one operation for each possible arity.
- This functor \mathbb{T} should be a monad.
- Selecting basic arities \Leftrightarrow picking a subfunctor $B \subseteq \mathbb{T}$.
- Generating $T \approx \bigcup_n (id \cup B)^n$, the smallest submonad of \mathbb{T} containing B .

Shapely functors: strategy

Find a subcategory \mathcal{C} of $\text{Endo}(\text{MGph})$

- stable under composition and
- having a **terminal** object \mathbb{T} .

I.e., such that $\forall C \in \mathcal{C}, \exists!$ morphism $C \rightarrow \mathbb{T}$.

Indeed:

- \mathbb{T} automatically a monad *via* $\mathbb{T} \circ \mathbb{T} \rightarrow \mathbb{T}$;
- can then generate $\bigcup_n B^n$ amongst subfunctors of \mathbb{T} .

Towards shapely functors I: analytic functors

Subcategory $\text{Analytic}(\text{MGph}) \subseteq \text{Endo}(\text{MGph})$ of functors s.t.

$$T(Y)_n = \sum_{x \in T(1)_n} \text{MGph}(A(x), Y)/G(x)$$

where

- $A(x)$ is the arity of x ,
- $G(x) \triangleleft \mathfrak{S}_{A(x)}$ is a subgroup of the automorphism group of $A(x)$.
- Generalisation of Joyal's analytic endofunctors on sets.

Miss again!

- Does have a terminal object.
- Not stable under composition.

Towards shapely functors II: cellular functors

Subcategory $\text{Cell}(\text{MGph}) \subseteq \text{Analytic}(\text{MGph}) \subseteq \text{Endo}(\text{MGph})$.

Miss again!

- Stable under composition.
- No terminal object!