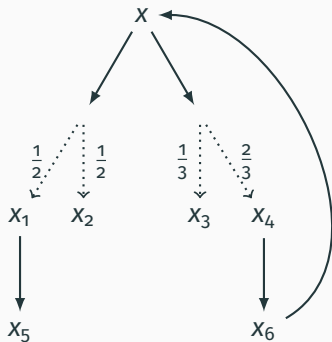


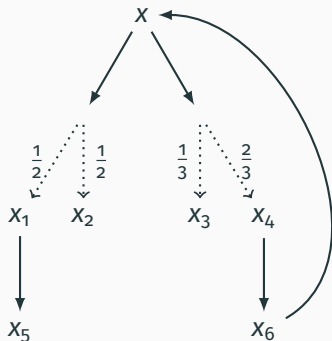
MONADS, EQUATIONAL THEORIES AND METRICS FOR NONDETERMINISTIC AND PROBABILISTIC SYSTEMS

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CNRS, ENS Lyon

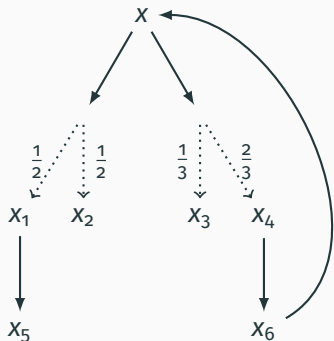
NONDETERMINISTIC AND PROBABILISTIC SYSTEMS





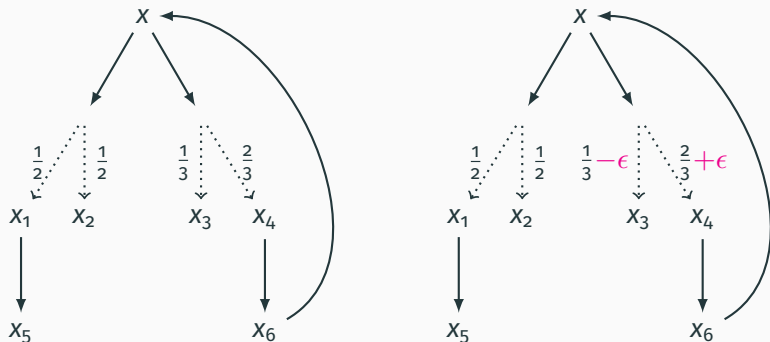
- nondeterminism and probability as computational effects:
monads and equational theories

NONDETERMINISTIC AND PROBABILISTIC SYSTEMS



- nondeterminism and probability as computational effects: monads and equational theories
- equational reasoning in program semantics

NONDETERMINISTIC AND PROBABILISTIC SYSTEMS



- nondeterminism and probability as computational effects:
monads and equational theories
- equational reasoning in program semantics
- program distance \Rightarrow monads on metric spaces and quantitative equational theories

**MONADS AND EQUATIONAL THEORIES FOR
NONDETERMINISM AND PROBABILITY**

Monad (\mathcal{M}, η, μ) in Sets

- functor $\mathcal{M} : X \mapsto \mathcal{M}(X)$
- unit $\eta_X : X \rightarrow \mathcal{M}(X)$
- multiplication $\mu_X : \mathcal{M}\mathcal{M}(X) \rightarrow \mathcal{M}(X)$

$$\begin{array}{ccccc}
 \mathcal{M}X & \xrightarrow{\eta\mathcal{M}} & \mathcal{M}^2X & \xleftarrow{\mathcal{M}\eta} & \mathcal{M}X \\
 & \parallel & \downarrow \mu & \parallel & \\
 & & \mathcal{M}X & &
 \end{array}$$

$$\begin{array}{ccc}
 \mathcal{M}^3X & \xrightarrow{\mu\mathcal{M}} & \mathcal{M}^2X \\
 \mathcal{M}\mu \downarrow & & \downarrow \mu \\
 \mathcal{M}^2X & \xrightarrow{\mu} & \mathcal{M}X
 \end{array}$$

Monad (\mathcal{M}, η, μ)
in Sets

Equational Theory (Σ, E)

- Σ a signature
- E a set of equations

- equations $t = s$
- deductive system: equational logic
 $\{t = s, s = u\} \vdash t = u$
- models: algebras (A, Σ^A) satisfying the equations

Monad (\mathcal{M}, η, μ)
in Sets



Equational Theory (Σ, E)

- Σ a signature
- E a set of equations

(Σ, E) is a presentation of (\mathcal{M}, η, μ)

The category $\mathbf{EM}(\mathcal{M})$ of Eilenberg-Moore algebras for (\mathcal{M}, η, μ) is isomorphic to the category $\mathbf{A}(\Sigma, E)$ of algebras (models) of (Σ, E)

Category $\mathbf{EM}(\mathcal{M})$

- objects: $(A, \alpha : \mathcal{M}(A) \rightarrow A)$
with α commuting with η, μ
- arrows: algebra morphisms

Category $\mathbf{A}(\Sigma, E)$

- objects: models (A, Σ^A) of (Σ, E)
- arrows: homomorphisms of (Σ, E) -algebras

MONADS AND EQUATIONAL THEORIES

Monad (\mathcal{M}, η, μ)
in Sets



Equational Theory (Σ, E)

- Σ a signature
- E a set of equations

(Σ, E) is a **presentation** of (\mathcal{M}, η, μ)

The category $\mathbf{EM}(\mathcal{M})$ of Eilenberg-Moore algebras for (\mathcal{M}, η, μ) is isomorphic to the category $\mathbf{A}(\Sigma, E)$ of algebras (models) of (Σ, E)

Corollary:

$$\mathcal{M}(X) \cong \text{Terms}(X, \Sigma) /_E$$

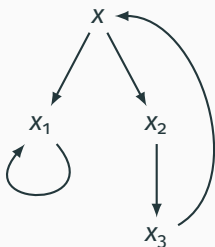
EXAMPLE: NONDETERMINISM

Monad (\mathcal{M}, η, μ)
in Sets



Equational Theory (Σ, E)

- Σ a signature
- E a set of equations



$$c : X \rightarrow \mathcal{P}(X)$$

$$c(x) = \{x_1, x_2\}$$

$$c(x_1) = \{x_1\}$$

...

EXAMPLE: NONDETERMINISM

Monad (\mathcal{M}, η, μ)
in Sets



Equational Theory (Σ, E)

- Σ a signature
- E a set of equations

Powerset (non-empty)
monad (\mathcal{P}, η, μ)

- $\mathcal{P} : X \mapsto \{S \mid S \text{ is a non-empty, finite subset of } X\}$
- $\eta : x \mapsto \{x\}$
- $\mu : \{S_1, \dots, S_n\} \mapsto \bigcup_i S_i$



Equational theory of semilattices

- $\Sigma =$ binary operation \oplus
- axioms of $E =$

$$\begin{array}{lcl} (x \oplus y) \oplus z & \stackrel{(A)}{=} & x \oplus (y \oplus z) \\ x \oplus y & \stackrel{(C)}{=} & y \oplus x \\ x \oplus x & \stackrel{(I)}{=} & x \end{array}$$

EXAMPLE: NONDETERMINISM

Monad (\mathcal{M}, η, μ)
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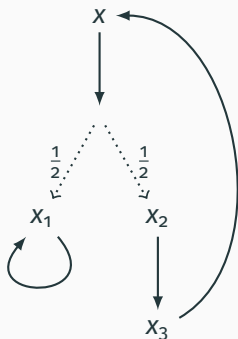
EXAMPLE: PROBABILITY

Monad (\mathcal{M}, η, μ)
in Sets



Equational Theory (Σ, E)

- Σ a signature
- E a set of equations



$$c : X \rightarrow \mathcal{D}(X)$$

$$c(x) = \frac{1}{2}x_1 + \frac{1}{2}x_2$$

$$c(x_1) = 1x_1$$

...

EXAMPLE: PROBABILITY

Monad (\mathcal{M}, η, μ)
in Sets



Equational Theory (Σ, E)

- Σ a signature
- E a set of equations

Distribution monad (\mathcal{D}, η, μ)

- $\mathcal{D} : X \mapsto \{\Delta \mid \Delta \text{ is a finitely supported probability distribution on } X\}$



Equational theory of convex algebras

- $\eta : X \mapsto 1X$
- $\mu : \sum_i p_i \Delta_i \mapsto \sum_i p_i \cdot \Delta_i$

- $\Sigma =$ binary operations $+_p$ for all $p \in (0, 1)$

- axioms of $E =$

$$(x +_q y) +_p z \stackrel{(A_p)}{=} x +_{pq} (y +_{\frac{p(1-q)}{1-pq}} z)$$
$$x +_p y \stackrel{(C_p)}{=} y +_{1-p} x$$
$$x +_p x \stackrel{(I_p)}{=} x$$

EXAMPLE: PROBABILITY

Monad (\mathcal{M}, η, μ)
in Sets



Equational Theory (Σ, E)

- Σ a signature
- E a set of equations

nondeterminism

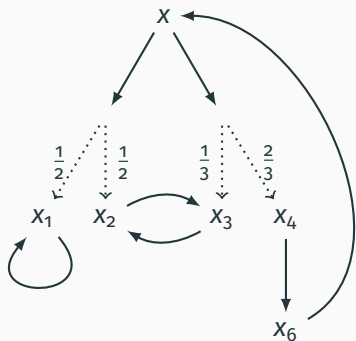
+

probability

=

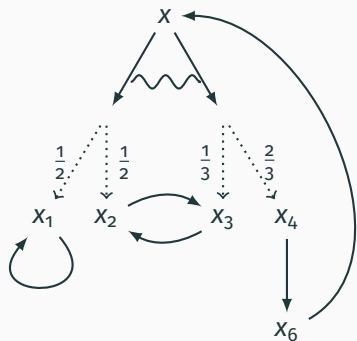
?

COMBINING NONDETERMINISM AND PROBABILITY



- a transition reaches a set of probability distributions $\{ \frac{1}{2}X_1 + \frac{1}{2}X_2, \frac{1}{3}X_3 + \frac{2}{3}X_4 \}$
- Problem: $\mathcal{P} \circ \mathcal{D}$ is not a monad [Varacca, Winskel 2006]

COMBINING NONDETERMINISM AND PROBABILITY

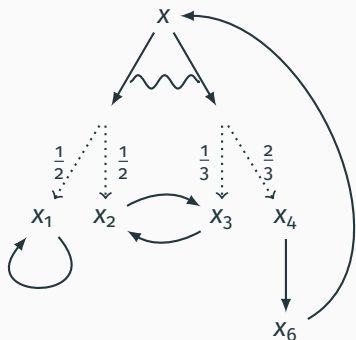


- a transition reaches a set of probability distributions $\left\{ \frac{1}{2}X_1 + \frac{1}{2}X_2, \frac{1}{3}X_3 + \frac{2}{3}X_4 \right\}$
- Problem: $\mathcal{P} \circ \mathcal{D}$ is not a monad [Varacca, Winskel 2006]

Solution: use **convex sets of probability distributions**

$$\left\{ \frac{1}{2}X_1 + \frac{1}{2}X_2, \dots, \frac{1}{4}X_1 + \frac{1}{4}X_2 + \frac{1}{6}X_3 + \frac{1}{3}X_4, \dots, \frac{1}{3}X_3 + \frac{2}{3}X_4 \right\}$$

COMBINING NONDETERMINISM AND PROBABILITY



- a transition reaches a set of probability distributions $\{ \frac{1}{2}X_1 + \frac{1}{2}X_2, \frac{1}{3}X_3 + \frac{2}{3}X_4 \}$
- Problem: $\mathcal{P} \circ \mathcal{D}$ is not a monad [Varacca, Winskel 2006]

Solution: use **convex sets of probability distributions**

$$\{ \frac{1}{2}X_1 + \frac{1}{2}X_2, \dots, \frac{1}{4}X_1 + \frac{1}{4}X_2 + \frac{1}{6}X_3 + \frac{1}{3}X_4, \dots, \frac{1}{3}X_3 + \frac{2}{3}X_4 \}$$

+ accounts for probabilistic schedulers

The monad (\mathcal{C}, η, μ) in Sets:

- $\mathcal{C} : X \mapsto \{S \mid S \text{ is a non-empty, convex-closed, finitely generated set of finitely supported probability distributions over } X\}$

- $\eta_X : X \rightarrow \mathcal{C}(X)$

$$\eta_X : x \mapsto \{ \mathbf{1}_x \}$$

- $\mu_X : \mathcal{C}\mathcal{C}(X) \rightarrow \mathcal{C}(X)$

$$\mu_X : \bigcup_i \{\Delta_i\} \mapsto \bigcup_i \text{WMS}(\Delta_i)$$

with $\text{WMS} : \mathcal{DC}(X) \rightarrow \mathcal{C}(X)$ the *weighted Minkowski sum*

$$\text{WMS}\left(\sum_{i=1}^n p_i S_i\right) = \left\{ \sum_{i=1}^n p_i \cdot \Delta_i \mid \text{for each } 1 \leq i \leq n, \Delta_i \in S_i \right\}$$

THE EQUATIONAL THEORY FOR NONDETERMINISM AND PROBABILITY

Monad (\mathcal{M}, η, μ)
in Sets



Equational Theory (Σ, E)

- Σ a signature
- E a set of equations

Convex sets (non-empty)
of distributions
monad (\mathcal{C}, η, μ)

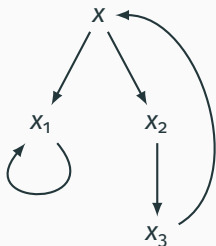


Equational theory of convex semilattices

- $\Sigma = \oplus$ and $+_p$ for all $p \in (0, 1)$
- axioms E :
 - axioms of semilattices
 - axioms of convex algebras
 - distributivity axiom (D)
$$(x \oplus y) +_p z \stackrel{(D)}{=} (x +_p z) \oplus (y +_p z)$$

[Bonchi, Sokolova, V. 2019 and 2020]

TERMINATION AND EQUATIONAL REASONING IN PROGRAM SEMANTICS



$$c : X \rightarrow \mathcal{P}(X)$$

$$c(x) = \{x_1, x_2\}$$

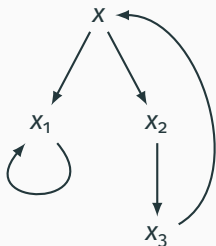
$$c(x_1) = \{x_1\}$$

...

Bisimulation equivalence (= coalgebraic behavioral equivalence)

$\exists R$ such that xRy implies:

- if $x \rightarrow x'$ then $y \rightarrow y'$ and $x'Ry'$
- symmetric



$$c : X \rightarrow \mathcal{P}(X)$$

$$c(x) = \{x_1, x_2\}$$

$$c(x_1) = \{x_1\}$$

...

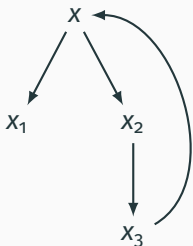
Bisimulation equivalence (= coalgebraic behavioral equivalence)

$\exists R$ such that xRy implies:

- if $x \rightarrow x'$ then $y \rightarrow y'$ and $x'Ry'$
- symmetric

trivial if all states can move \Rightarrow we need some observable: termination

NONDETERMINISTIC SYSTEMS AND BISIMULATION



$$c : X \rightarrow \mathcal{P}^{\emptyset}(X)$$

$$c(x) = \{x_1, x_2\}$$

$$c(x_1) = \emptyset$$

...

i.e., $x_1 \not\rightarrow$

Bisimulation equivalence (= coalgebraic behavioral equivalence)

$\exists R$ such that xRy implies:

■ if $x \rightarrow x'$ then $y \rightarrow y'$ and $x'Ry'$

■ symmetric

trivial if all states can move \Rightarrow we need some observable: termination

NONDETERMINISM + TERMINATION

Monad (\mathcal{M}, η, μ)
in Sets



Equational Theory (Σ, E)

- Σ a signature
- E a set of equations

Powerset (possibly empty)
monad $(\mathcal{P}^\emptyset, \eta, \mu)$

- $\mathcal{P}^\emptyset : X \mapsto \{S \mid S \text{ is a finite subset of } X\}$
- $\eta : x \mapsto \{x\}$
- $\mu : \{S_1, \dots, S_n\} \mapsto \bigcup_i S_i$



Equational theory of
semilattices with bottom

- $\Sigma = \star, \oplus$
- axioms of $E =$
 - axioms of semilattices

$$\begin{array}{ccc} (x \oplus y) \oplus z & \stackrel{(A)}{=} & x \oplus (y \oplus z) \\ x \oplus y & \stackrel{(C)}{=} & y \oplus x \\ x \oplus x & \stackrel{(I)}{=} & x \end{array}$$

- bottom axiom $x \oplus \star = x$

Monad (\mathcal{M}, η, μ)
in Sets



Equational Theory (Σ, E)

- Σ a signature
- E a set of equations

Convex sets
(possibly empty)
of distributions
monad $(\mathcal{C}^\emptyset, \eta, \mu)$



Equational theory of convex semilattices
with bottom and black-hole

- $\Sigma = \star, \oplus, +_p$ for all $p \in (0, 1)$
- axioms of $E =$
 - axioms of convex semilattices
 - bottom axiom $x \oplus \star = x$
 - black-hole axiom $x +_p \star = \star$

[Mio, Sarkis, V. 2021]

APPLICATION: EQUATIONAL REASONING ON PROCESSES

Processes $P ::= \mathbf{nil} \mid a.P \mid P_1 \bar{\oplus} P_2 \mid P_1 \bar{+}_p P_2$

Semantics $c : Proc \rightarrow Terms(Proc, \Sigma)$

with $\Sigma = \star, \oplus, +_p$

$$\mathbf{nil} \frac{}{\mathbf{nil} \rightarrow \star}$$

$$\bar{\oplus} \frac{P_1 \rightarrow t_1 \quad P_2 \rightarrow t_2}{P_1 \bar{\oplus} P_2 \rightarrow t_1 \oplus t_2}$$

$$\text{action} \frac{}{a.P \rightarrow P}$$

$$\bar{+}_p \frac{P_1 \rightarrow t_1 \quad P_2 \rightarrow t_2}{P_1 \bar{+}_p P_2 \rightarrow t_1 +_p t_2}$$

APPLICATION: EQUATIONAL REASONING ON PROCESSES

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$$(a.(nil \bar{\oplus} nil) \bar{+}_{\frac{1}{3}} a^3.nil) \bar{\oplus} (a^2.nil \bar{+}_{\frac{1}{2}} a.nil)$$



$$((nil \bar{\oplus} nil) +_{\frac{1}{3}} a^2.nil) \bar{\oplus} (a.nil +_{\frac{1}{2}} nil)$$

APPLICATION: EQUATIONAL REASONING ON PROCESSES

Processes $P ::= \mathbf{nil} \mid a.P \mid P_1 \bar{\oplus} P_2 \mid P_1 \bar{+}_p P_2$

Semantics $c : Proc \rightarrow Terms(Proc, \Sigma)_{/E}$

with $\Sigma = \star, \oplus, +_p$

$(a.(\mathbf{nil} \bar{\oplus} \mathbf{nil}) \bar{+}_{\frac{1}{3}} a^3.\mathbf{nil}) \bar{\oplus} (a^2.\mathbf{nil} \bar{+}_{\frac{1}{2}} a.\mathbf{nil})$



$[((\mathbf{nil} \bar{\oplus} \mathbf{nil}) +_{\frac{1}{3}} a^2.\mathbf{nil}) \bar{\oplus} (a.\mathbf{nil} +_{\frac{1}{2}} \mathbf{nil})]_{/E}$

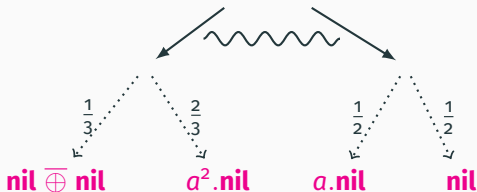
APPLICATION: EQUATIONAL REASONING ON PROCESSES

Processes $P ::= \mathbf{nil} \mid a.P \mid P_1 \bar{\oplus} P_2 \mid P_1 \bar{+}_p P_2$

Semantics $c : Proc \rightarrow Terms(Proc, \Sigma) / E \simeq \mathcal{C}^\emptyset(Proc)$

with $\Sigma = \star, \oplus, +_p$

$(a.(nil \bar{\oplus} nil) \bar{+}_{\frac{1}{3}} a^3.nil) \bar{\oplus} (a^2.nil \bar{+}_{\frac{1}{2}} a.nil)$



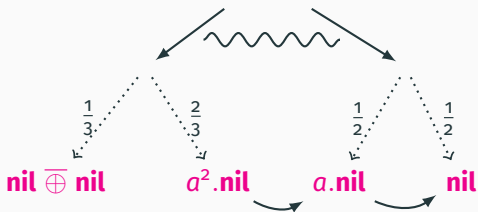
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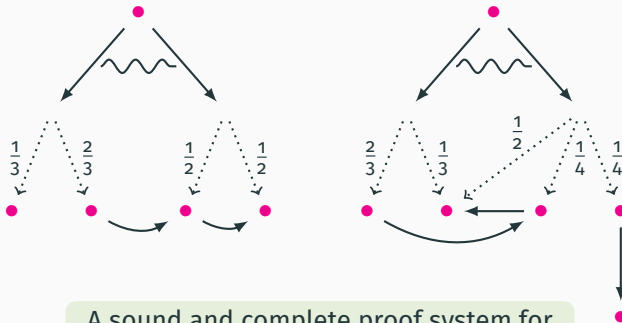


APPLICATION: EQUATIONAL REASONING ON PROCESSES

Processes $P ::= \mathbf{nil} \mid a.P \mid P_1 \oplus P_2 \mid P_1 \overline{+}_p P_2$

Semantics $c : Proc \rightarrow Terms(Proc, \Sigma) / E \simeq \mathcal{C}^\emptyset(Proc)$

with $\Sigma = \star, \oplus, +_p$



A sound and complete proof system for bisimulation equivalence on processes

$$P \sim P' \text{ iff } P \stackrel{E}{\equiv} P'$$

AXIOMS FOR TERMINATION: THE ISSUE WITH BLACK-HOLE

Processes $P ::= \mathbf{nil} \mid a.P \mid P_1 \bar{\oplus} P_2 \mid P_1 \bar{+}_p P_2$

Semantics $c : Proc \rightarrow Terms(Proc, \Sigma) / E \simeq C^\emptyset(Proc)$

with $\Sigma = \star, \oplus, +_p$

Bottom axiom $x \oplus \star = x$

Black-hole axiom $x +_p \star = \star$

AXIOMS FOR TERMINATION: THE ISSUE WITH BLACK-HOLE

Processes $P ::= \mathbf{nil} \mid a.P \mid P_1 \oplus P_2 \mid P_1 \overline{+}_p P_2$

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Bottom axiom $x \oplus \star = x$

Black-hole axiom $x +_p \star = \star$

$a.\mathbf{nil} \overline{+}_{\frac{1}{2}} \mathbf{nil}$



$[\mathbf{nil} +_{\frac{1}{2}} \star]_{/E}$

\mathbf{nil}



$[\star]_{/E}$

AXIOMS FOR TERMINATION: THE ISSUE WITH BLACK-HOLE

Processes $P ::= \mathbf{nil} \mid a.P \mid P_1 \bar{\oplus} P_2 \mid P_1 \bar{+}_p P_2$

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Bottom axiom $x \oplus \star = x$

Black-hole axiom $x +_p \star = \star$

$$\begin{array}{ccc} a.\mathbf{nil} \bar{+}_{\frac{1}{2}} \mathbf{nil} & \sim & \mathbf{nil} \\ \downarrow & & \downarrow \\ [\mathbf{nil} +_{\frac{1}{2}} \star]_{/E} & = & [\star]_{/E} \end{array}$$

NONDETERMINISM + PROBABILITY + TERMINATION, BOTTOM ONLY

Monad (\mathcal{M}, η, μ)
in Sets



Equational Theory (Σ, E)

- Σ a signature
- E a set of equations

\perp -closed convex sets
(possibly empty)
of subdistributions
monad $(\mathcal{C}^\perp, \eta, \mu)$



Equational theory of convex semilattices
with bottom

- $\Sigma = \star, \oplus, +_p$ for all $p \in (0, 1)$
- axioms E :
 - axioms of convex semilattices
 - bottom axiom $x \oplus \star = x$

[Mio, Sarkis, V. 2021]

- subdistribution = $\sum_i p_i x_i$ with $\sum_i p_i \leq 1$
- S is \perp -closed = if $\sum_i p_i x_i \in S$ then $\sum_i q_i x_i \in S$ with $q_i \leq p_i$

BOTTOM ONLY SEMANTICS

Processes $P ::= \mathbf{nil} \mid a.P \mid P_1 \bar{\oplus} P_2 \mid P_1 \bar{+}_p P_2$

Semantics $c : Proc \rightarrow Terms(Proc, \Sigma)_{/E} \simeq C^\downarrow(Proc)$

with $\Sigma = \star, \oplus, +_p$

Bottom axiom $x \oplus \star = x$

BOTTOM ONLY SEMANTICS

Processes $P ::= \mathbf{nil} \mid a.P \mid P_1 \bar{\oplus} P_2 \mid P_1 \bar{+}_p P_2$

Semantics $c : Proc \rightarrow Terms(Proc, \Sigma)_{/E} \simeq \mathcal{C}^\downarrow(Proc)$

with $\Sigma = \star, \oplus, +_p$

Bottom axiom $x \oplus \star = x$

$$\begin{array}{ccc} a.\mathbf{nil} \bar{+}_{\frac{1}{2}} \mathbf{nil} & \not\sim & \mathbf{nil} \\ \downarrow & & \downarrow \\ [\mathbf{nil} +_{\frac{1}{2}} \star]_{/E} & \neq & [\star]_{/E} \end{array}$$

BOTTOM ONLY SEMANTICS

Processes $P ::= \mathbf{nil} \mid a.P \mid P_1 \bar{\oplus} P_2 \mid P_1 \bar{+}_p P_2$

Semantics $c : \mathit{Proc} \rightarrow \mathit{Terms}(\mathit{Proc}, \Sigma)_{/E} \simeq \mathcal{C}^\downarrow(\mathit{Proc})$

with $\Sigma = \star, \oplus, +_p$

Bottom axiom $x \oplus \star = x$

$a.\mathbf{nil} \bar{+}_{\frac{1}{2}} \mathbf{nil} \quad \not\sim \quad \mathbf{nil}$

↓

$\frac{1}{2} \downarrow \vdots \downarrow$
 \mathbf{nil}

Convex sets (non-empty)
of distributions
monad (\mathcal{C}, η, μ)



Equational theory of convex semilattices

Convex sets
(possibly empty)
of distributions
monad $(\mathcal{C}^\emptyset, \eta, \mu)$



Equational theory of convex semilattices
with bottom $x \oplus \star = x$
and black-hole $x +_p \star = \star$

\perp -closed convex sets
(possibly empty)
of subdistributions
monad $(\mathcal{C}^\downarrow, \eta, \mu)$



Equational theory of convex semilattices
with bottom $x \oplus \star = x$

MONADS ON METRIC SPACES AND PROGRAM DISTANCES

Category Met of Metric Spaces:

- objects (X, d)
- morphisms $f : (X, d) \rightarrow (Y, d')$ are non-expansive maps between metric spaces

$$f : X \rightarrow Y \quad \text{with} \quad d'(f(x_1), f(x_2)) \leq d(x_1, x_2)$$

Monad (\mathcal{M}, η, μ) in Sets lifted to a monad $(\hat{\mathcal{M}}, \hat{\eta}, \hat{\mu})$ in Met, with:

- $\hat{\mathcal{M}} : (X, d) \mapsto (\mathcal{M}(X), \text{lift}_{\mathcal{M}}(d))$
- $\hat{\eta}_{(X,d)} : (X, d) \rightarrow (\mathcal{M}(X), \text{lift}_{\mathcal{M}}(d))$ and
 $\hat{\mu}_{(X,d)} : (\mathcal{M}\mathcal{M}(X), \text{lift}_{\mathcal{M}}\text{lift}_{\mathcal{M}}(d)) \rightarrow (\mathcal{M}(X), \text{lift}_{\mathcal{M}}(d))$
non-expansive

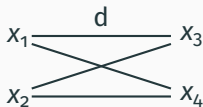
The monad (\mathcal{C}, η, μ) of convex sets of distributions can be lifted to a monad $(\hat{\mathcal{C}}, \hat{\eta}, \hat{\mu})$ in Met :

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THE MONAD OF CONVEX SETS OF DISTRIBUTIONS, ON METRIC SPACES

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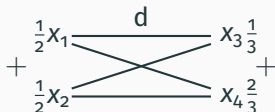
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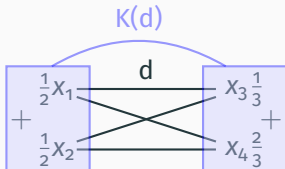
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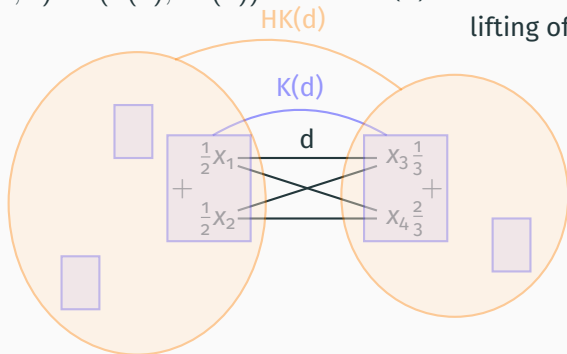
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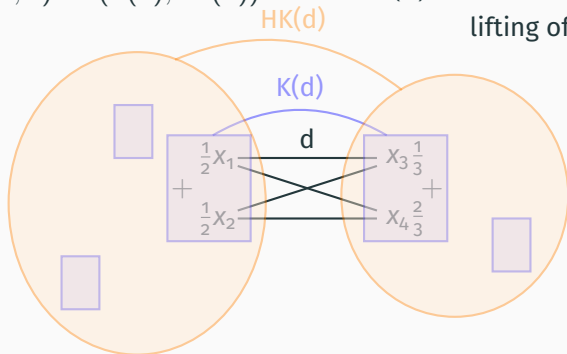
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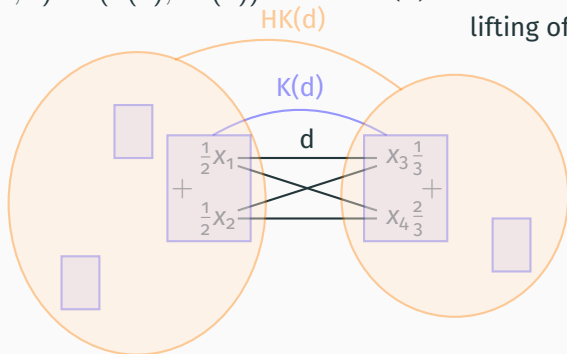


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Monad $(\hat{\mathcal{M}}, \hat{\eta}, \hat{\mu})$
in Met



Quantitative Equational Theory (Σ, Q)

- Σ a signature
- Q a set of quantitative inferences

Monad $(\hat{\mathcal{M}}, \hat{\eta}, \hat{\mu})$
in Met



Quantitative Equational Theory (Σ, Q)

- Σ a signature
- Q a set of quantitative inferences
- quantitative equations $t =_{\epsilon} t'$
- quantitative inferences

$$\{t_i =_{\epsilon_i} s_i\}_{i \in I} \vdash t =_{\epsilon} s$$
- quantitative deduction system
 - (Reflexivity) $\emptyset \vdash t =_0 t$
 - (Symmetry) $\{t =_{\epsilon} s\} \vdash s =_{\epsilon} t$
 - (Triangular) $\{t =_{\epsilon_1} u, u =_{\epsilon_2} s\} \vdash t =_{\epsilon_1 + \epsilon_2} s$
- models: quantitative algebras (A, Σ^A, d) satisfying the quantitative inferences

[Mardare, Panangaden, Plotkin 2016]

Monad $(\hat{\mathcal{M}}, \hat{\eta}, \hat{\mu})$
in \mathbf{Met}



Quantitative Equational Theory (Σ, Q)

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- Q a set of quantitative inferences

(Σ, Q) is a presentation of $(\hat{\mathcal{M}}, \hat{\eta}, \hat{\mu})$

The category $\mathbf{EM}(\hat{\mathcal{M}})$ of Eilenberg-Moore algebras for $(\hat{\mathcal{M}}, \hat{\eta}, \hat{\mu})$ is isomorphic to the category $\mathbf{A}(\Sigma, E)$ of quantitative algebras (models) of (Σ, Q)

THE QUANTITATIVE EQUATIONAL THEORY OF CONVEX SEMILATTICES

Monad $(\hat{\mathcal{M}}, \hat{\eta}, \hat{\mu})$
in Met



Quantitative Equational Theory (Σ, Q)

- Σ a signature
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Convex sets
(non-empty)
of distributions
monad $(\hat{\mathcal{C}}, \hat{\eta}, \hat{\mu})$
in Met



Quantitative equational theory
of convex semilattices

- $\Sigma = \oplus$ and $+_p$ for all $p \in (0, 1)$
- quantitative inferences $Q =$
 - axioms of convex semilattices,
with $t = t'$ becoming $\emptyset \vdash t =_o t'$
 - $\{X_1 =_{\epsilon_1} Y_1, X_2 =_{\epsilon_2} Y_2\} \vdash X_1 \oplus X_2 =_{\max(\epsilon_1, \epsilon_2)} Y_1 \oplus Y_2$
 - $\{X_1 =_{\epsilon_1} Y_1, X_2 =_{\epsilon_2} Y_2\} \vdash X_1 +_p X_2 =_{p \cdot \epsilon_1 + (1-p) \cdot \epsilon_2} Y_1 +_p Y_2$

[Mio, V. 2020]

THE QUANTITATIVE EQUATIONAL THEORY OF CONVEX SEMILATTICES

Monad $(\hat{\mathcal{M}}, \hat{\eta}, \hat{\mu})$
in Met



Quantitative Equational Theory (Σ, Q)

- Σ a signature
- Q a set of quantitative inferences

Corollary:

$$(\mathcal{C}(X), HK(d)) \cong (Terms(X, \Sigma) /_{E}, d_{(\Sigma, Q)})$$

with $d_{(\Sigma, Q)} = (t, t') \mapsto \inf\{\epsilon \mid \emptyset \vdash t =_{\epsilon} t'\}$

RECAP: ADDING TERMINATION, IN SETS

Monad (\mathcal{M}, η, μ)
in Sets



Equational Theory (Σ, E)

- Σ a signature
- E a set of equations

Convex sets
(possibly empty)
of distributions
monad $(\mathcal{C}^\emptyset, \eta, \mu)$



Equational theory of convex semilattices

with bottom $x \oplus \star = x$

and black-hole $x +_p \star = \star$

\perp -closed convex sets
(possibly empty)
of subdistributions
monad $(\mathcal{C}^\perp, \eta, \mu)$



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Negative results:

- The quantitative equational theory of convex semilattices with bottom and black-hole is trivial
- The multiplication μ of \mathcal{C}^\emptyset is not non-expansive \Rightarrow the same monad cannot be lifted to Met

LIFTING TO MET: NEGATIVE RESULTS

Convex sets
(possibly empty)
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Equational theory of convex semilattices
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Monad $(\hat{\mathcal{M}}, \hat{\eta}, \hat{\mu})$
in Met



Quantitative Equational Theory (Σ, Q)

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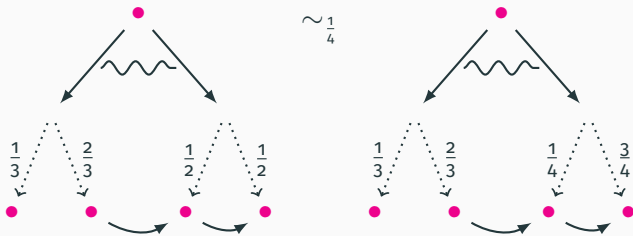
Quantitative equational theory of convex
semilattices with bottom $x \oplus \star = x$

APPLICATION: EQUATIONAL REASONING FOR BISIMULATION METRIC

Processes $P ::= \mathbf{nil} \mid a.P \mid P_1 \oplus P_2 \mid P_1 \overline{+}_p P_2$

Semantics $c : Proc \rightarrow Terms(Proc, \Sigma) / Q$

with $\Sigma = \star, \oplus, +_p$



Convex sets (non-empty)
of distributions
monad (\mathcal{C}, η, μ)



Equational theory of convex semilattices

Convex sets
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Equational theory of convex semilattices
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RECAP

YES in Met

Convex sets (non-empty)
of distributions
monad (\mathcal{C}, η, μ)



Equational theory of convex semilattices

NO in Met

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Equational theory of convex semilattices
with bottom $x \oplus \star = x$
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YES in Met

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Equational theory of convex semilattices
with bottom $x \oplus \star = x$

- more process operators, equivalences, metrics
- equational reasoning for trace equivalences and metrics
- recover convex sets of distributions monad in Met, and its presentation, compositionally [Goy, Petrisan 2020]

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Thank you!